

# Digital Communication Systems

## ECS 452

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**Source Coding**



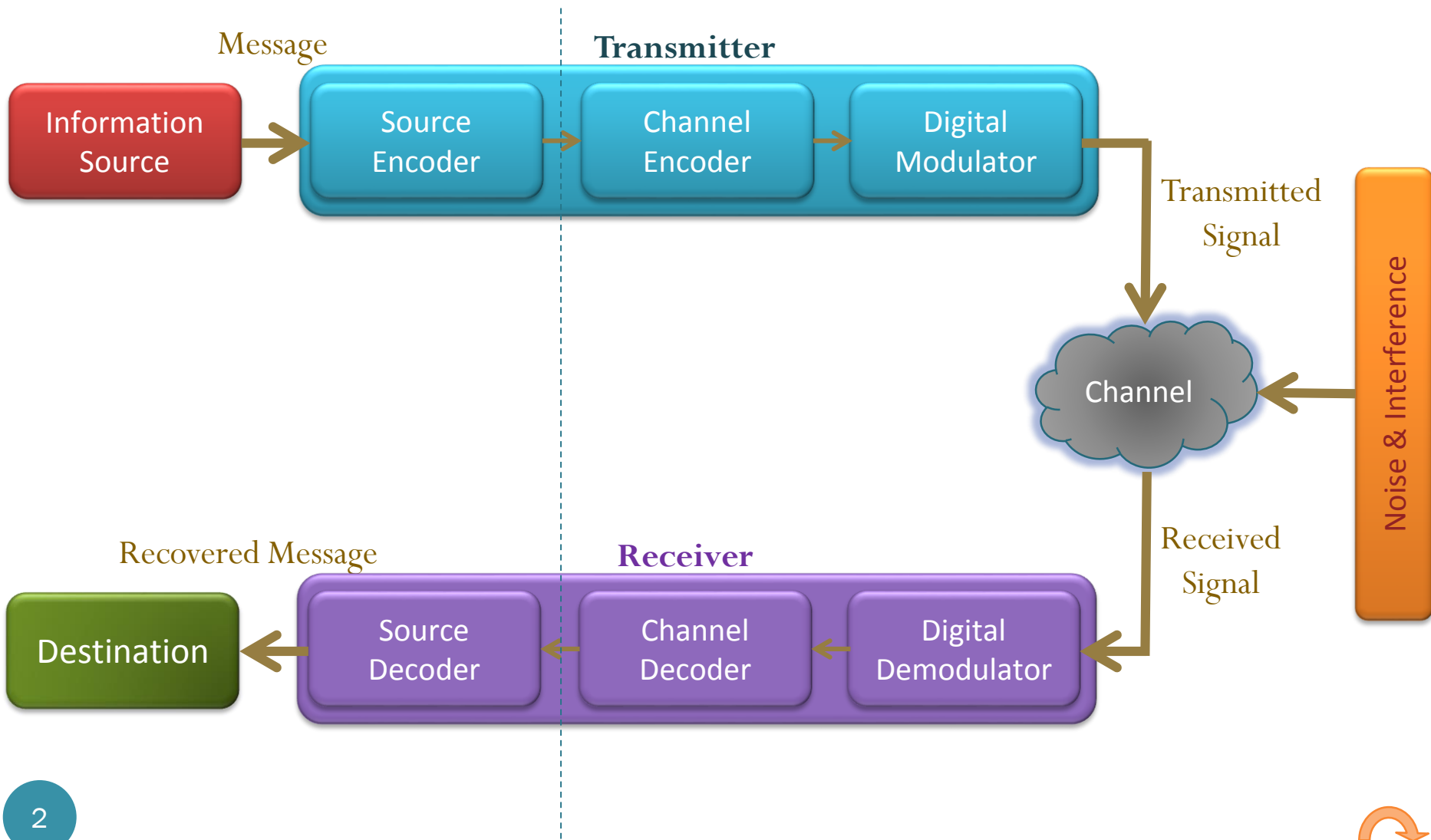
**Office Hours:**

**BKD 3601-7**

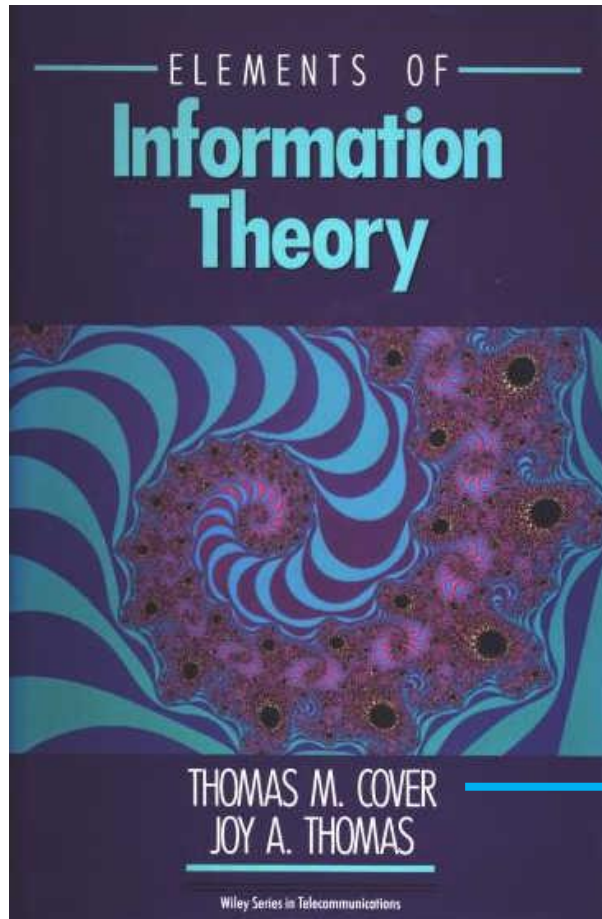
**Monday 14:00-16:00**

**Wednesday 14:40-16:00**

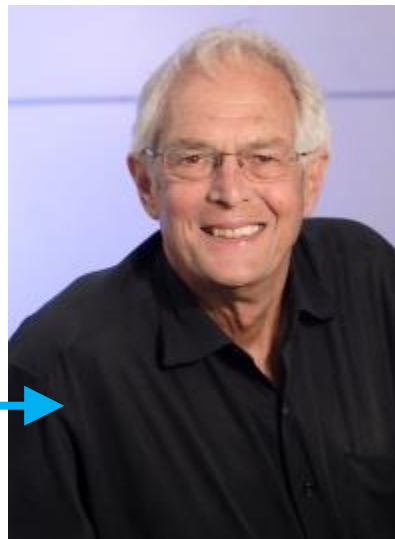
# Elements of digital commu. sys.



# Reference



- Elements of Information Theory
- 2006, 2nd Edition
- Chapters 2, 4 and 5



‘the jewel in Stanford's crown’

One of the greatest information theorists since Claude Shannon (and the one most like Shannon in approach, clarity, and taste).





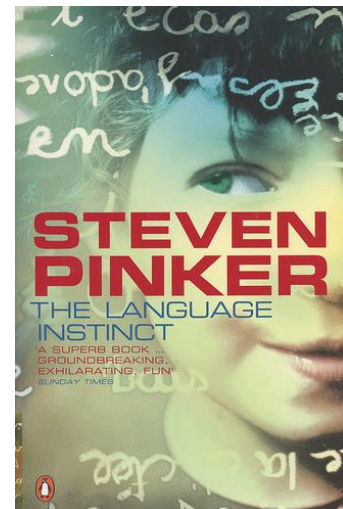
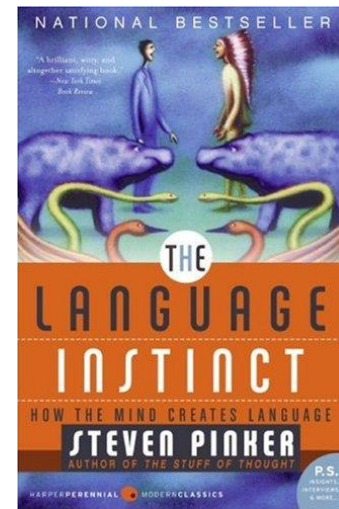
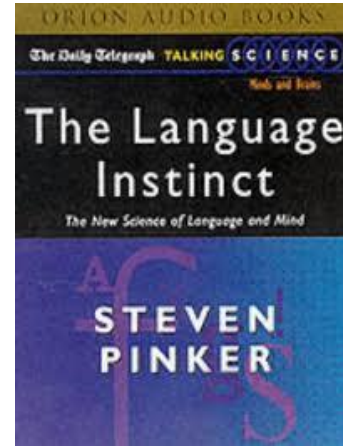
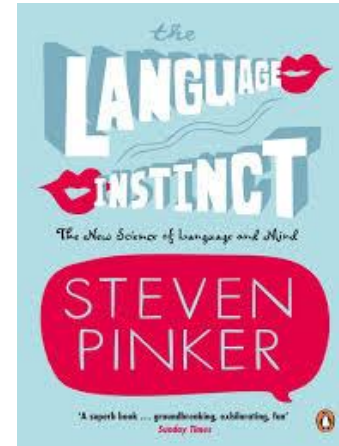
# English Redundancy: Ex. 1

J-st tr- t- r--d th-s s-nt-nc-.



# English Redundancy: Ex. 2

yxx cxn xndxrstxnd  
whxt x xm wrxtxng  
xvxn xf x rxplxcx xll  
thx vxwxls wxth xn 'x'  
(t gts lttl hrdr f y dn't  
vn kn whr th vwls r).



# English Redundancy: Ex. 3

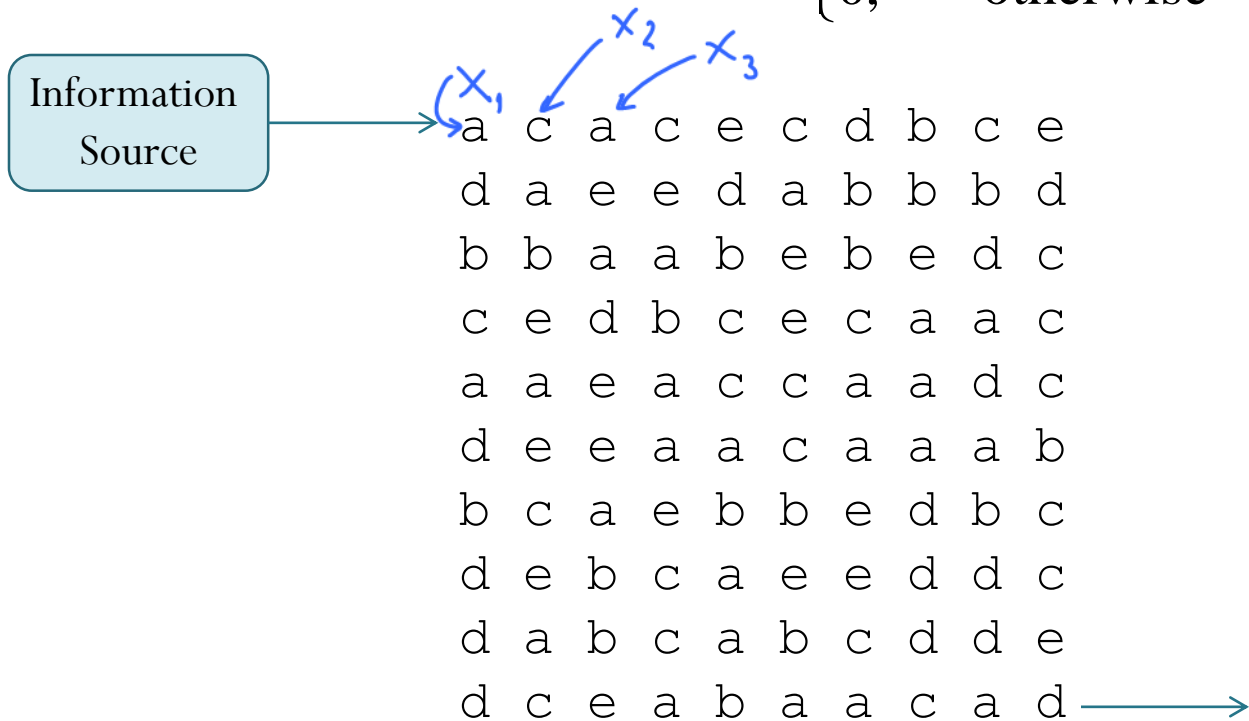
To be, or xxx xx xx,  
xxxx xx xxx xxxxxxxx



# Ex. DMS (1)

$$\mathcal{S}_X = \{a, b, c, d, e\}$$

$$p_X(x) = \begin{cases} 1/5, & x \in \{a, b, c, d, e\} \\ 0, & \text{otherwise} \end{cases}$$



Approximately 20% are letter 'a's





# Ex. DMS (2)

$$\mathcal{S}_X = \{1, 2, 3, 4\}$$

$$p_X(x) = \begin{cases} 1/2, & x = 1, \\ 1/4, & x = 2, \\ 1/8, & x \in \{3, 4\} \\ 0, & \text{otherwise} \end{cases}$$

Information Source



2 1 1 2 1 4 1 1 1 1  
 1 1 4 1 1 2 4 2 2 1  
 3 1 1 2 3 2 4 1 2 4  
 2 1 1 2 1 1 3 3 1 1  
 1 3 4 1 4 1 1 2 4 1  
 4 1 4 1 2 2 1 4 2 1  
 4 1 1 1 1 2 1 4 2 4  
 2 1 1 1 2 1 2 1 3 2  
 2 1 1 1 1 1 1 2 3 2  
 2 1 1 2 1 4 2 1 2 1



Approximately 50% are number '1's



# Demo: DMS in MATLAB

```
clear all; close all;
```

```
S_X = [1 2 3 4]; p_X = [1/2 1/4 1/8 1/8]; n = 1e6;
```

```
SourceString = randsrc(1,n,[S_X;p_X]);
```

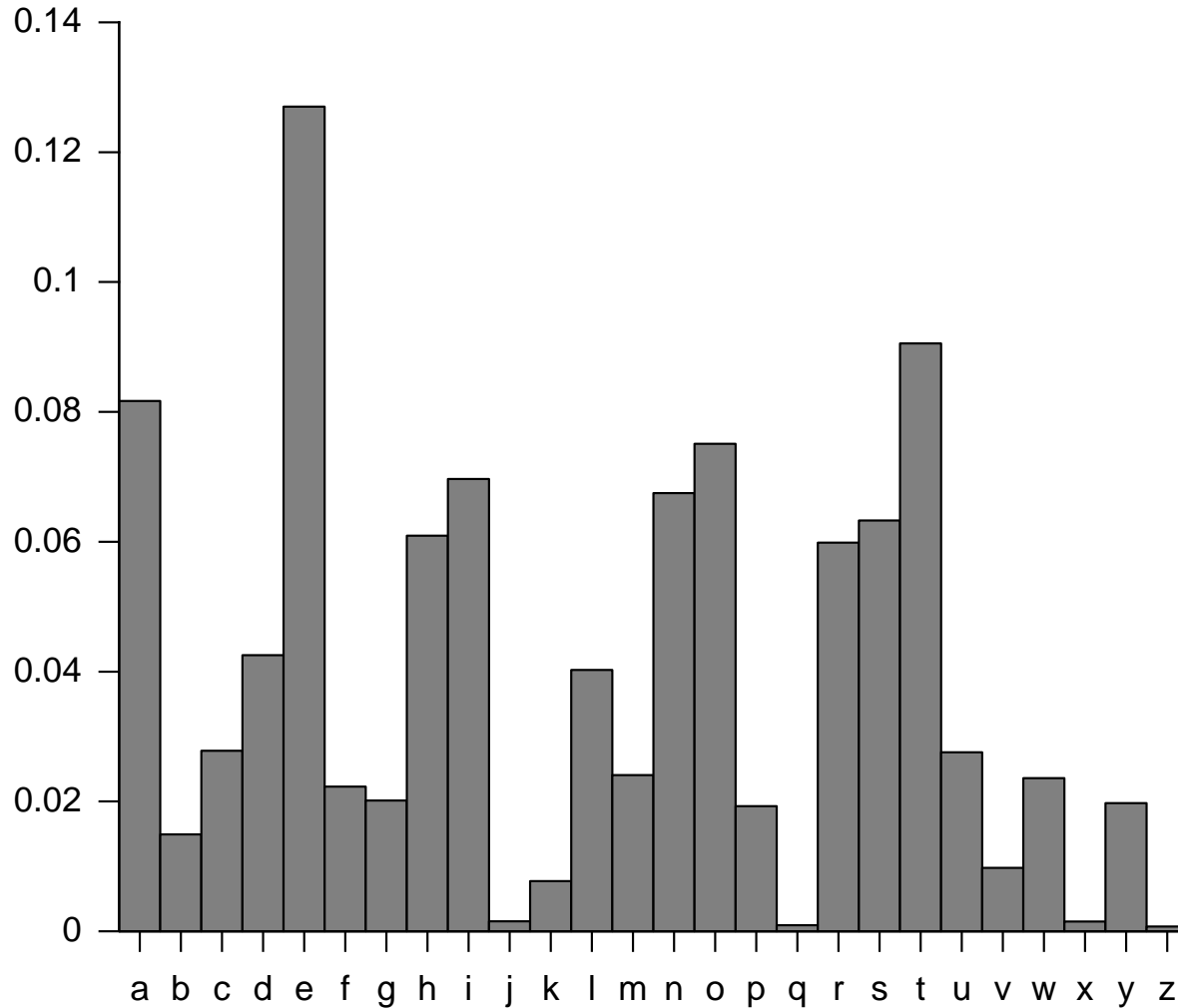
Alternatively, we can also use

```
SourceString = datasample(S_X,n,'Weights',p_X);
```

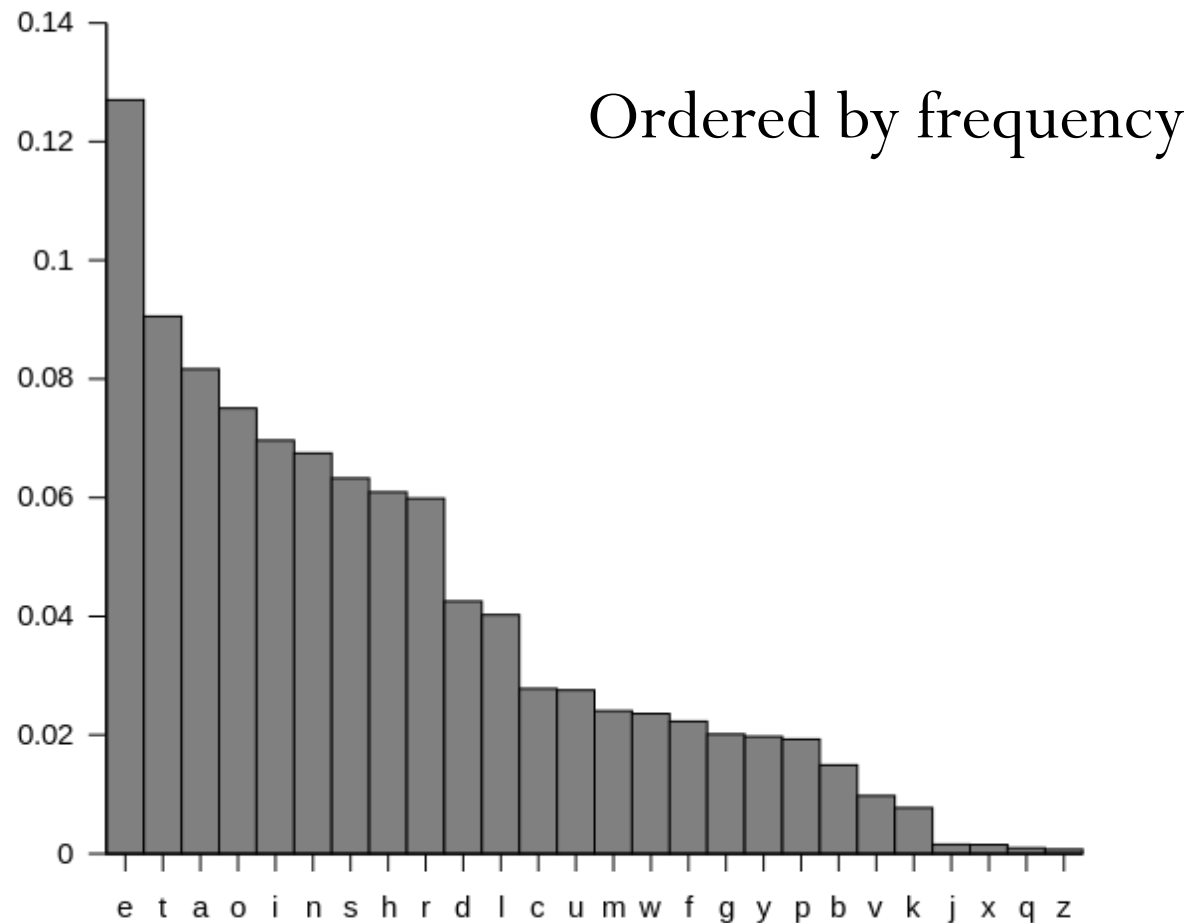
```
rf = hist(SourceString,S_X)/n; % Ref. Freq. calc.  
stem(S_X,rf,'rx','LineWidth',2) % Plot Rel. Freq.  
hold on  
stem(S_X,p_X,'bo','LineWidth',2) % Plot pmf  
xlim([min(S_X)-1,max(S_X)+1])  
legend('Rel. freq. from sim.','pmf p_X(x)')  
xlabel('x')  
grid on
```



# Relative freq. of letters in the English language



# Relative freq. of letters in the English language

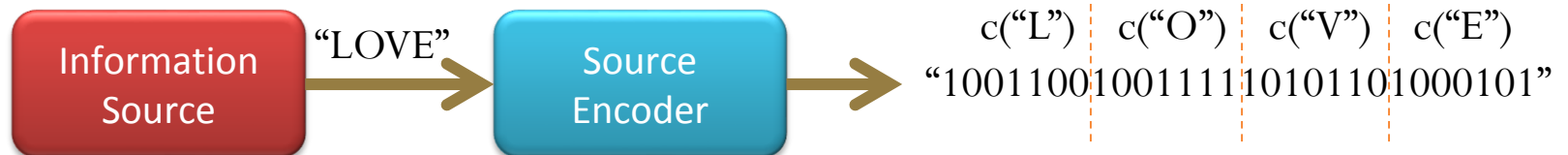


# Example: ASCII Encoder

Character	Codeword
:	
E	1000101
:	
L	1001100
:	
O	1001111
:	
V	1010110
:	

MATLAB:

```
>> M = 'LOVE';  
>> X = dec2bin(M, 7);  
>> X = reshape(X', 1, numel(X))  
X =  
1001100100111110101101000101
```



# Morse code

(wired and wireless)

- **Telegraph network**
- Samuel **Morse**, 1838
- A sequence of on-off tones (or , lights, or clicks)



A	● —	U	● ● —
B	— ● ● ●	V	● ● ● —
C	— — ● — ●	W	● — — —
D	— — ● ●	X	— ● ● — —
E	●	Y	— — — — ●
F	● ● — — ●	Z	— — — ● ●
G	— — — ●		
H	● ● ● ●		
I	● ●		
J	● — — — — —		
K	— ● — — —	1	● — — — — —
L	● — — ● ●	2	● ● — — — —
M	— — — —	3	● ● ● — — —
N	— — ●	4	● ● ● ● — —
O	— — — — —	5	● ● ● ● ●
P	● — — — — ●	6	— — ● ● ● ●
Q	— — — ● — —	7	— — — ● ● ●
R	● — — ● ●	8	— — — — ● ●
S	● ● ● ●	9	— — — — — ●
T	— — — — —	0	— — — — — —



# Example



"I love you." in Morse code



Examples Random

Input interpretation:

Morse code I love you.

Morse code translation:

•• | | •—•• | — — | ••• — | • | | —• — | — — |  
I | L | O | V | E | Y | O |  
•• — | • — • — • —  
U | .

Download page

POWERED BY THE WOLFRAM LANGUAGE



# Morse code: Key Idea

Frequently-used characters (e,t) are mapped to short codewords.

A	• —	U	• • —
B	— • • •	V	• • • —
C	— • — •	W	• — —
D	— • •	X	— • • —
E	•	Y	— • — —
F	• • — •	Z	— — • •
G	— — •		
H	• • • •		
I	• •		
J	• — — —		
K	— • — —		
L	• — • •		
M	— —		
N	— •		
O	— — —		
P	• — — •		
Q	— — • —		
R	• — •		
S	• • •		
T	—		
		1	• — — — —
		2	• • — — —
		3	• • • — —
		4	• • • • —
		5	• • • • •
		6	— • • • •
		7	— — • • •
		8	— — — • •
		9	— — — — •
		0	— — — — —

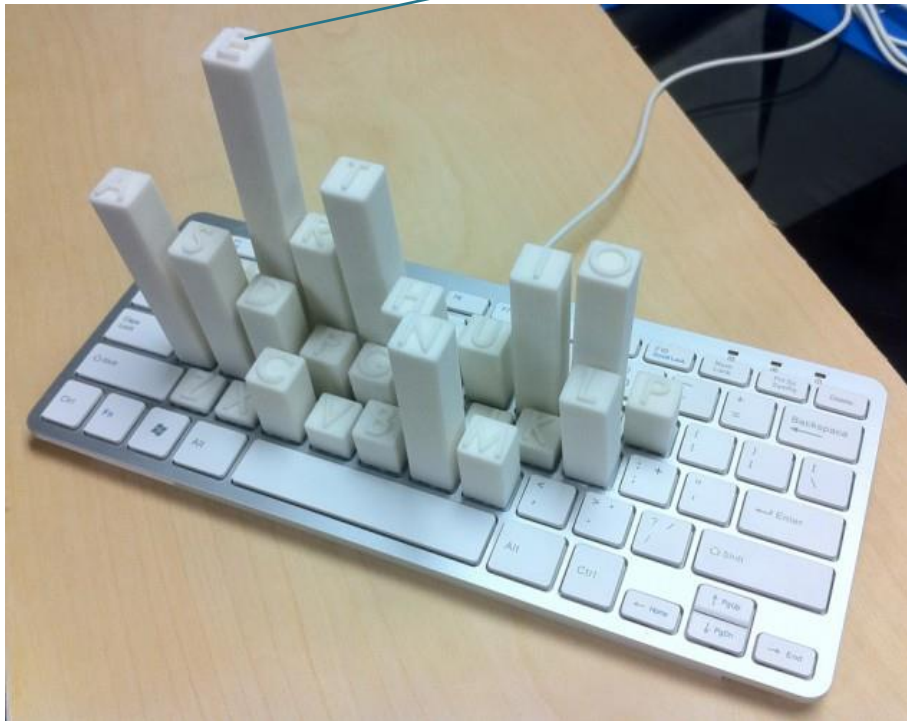
Basic form of compression.





# Morse code: Key Idea

Frequently-used characters are mapped to short codewords.

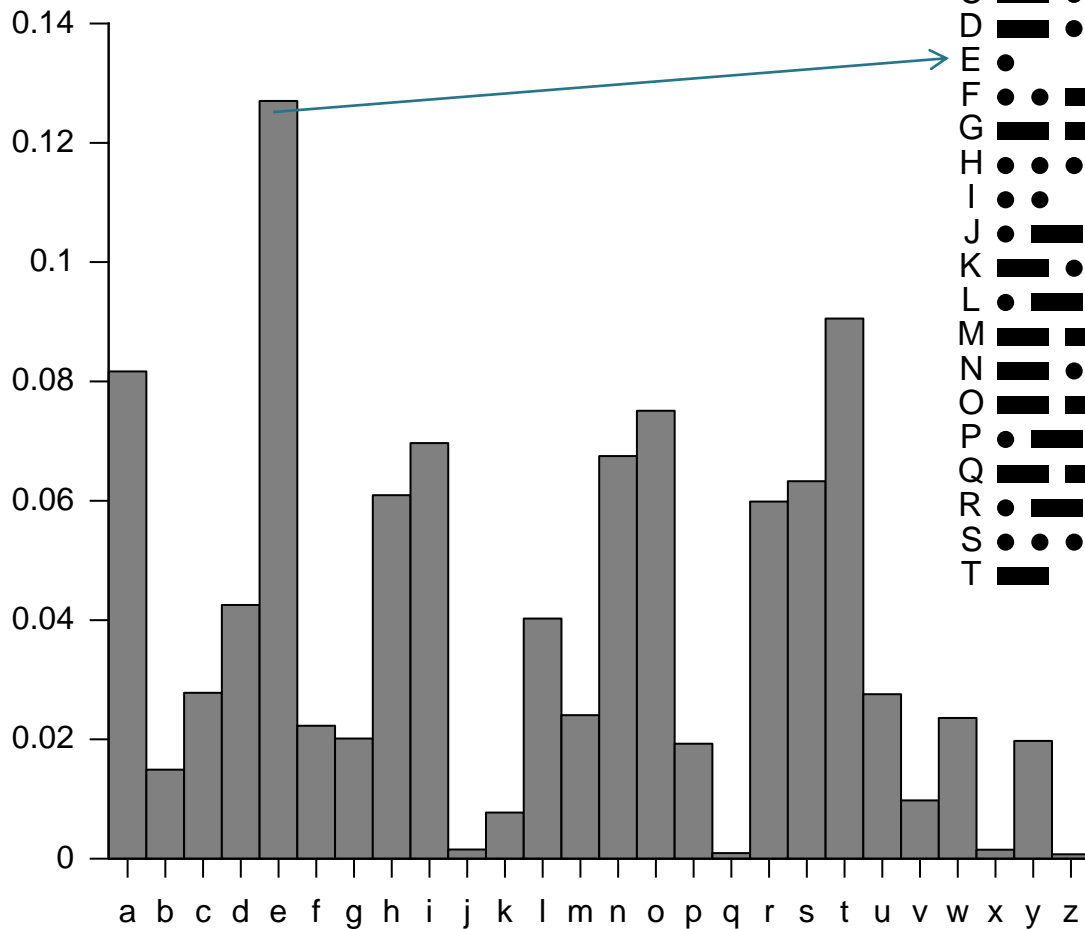


A	● —	U	● ● —
B	— ● ● ●	V	● ● ● —
C	— ● — ●	W	● — —
D	— ● ●	X	— ● ● —
E	●	Y	— ● — —
F	● ● — ●	Z	— — ● ●
G	— — ●		
H	● ● ● ●		
I	● ●		
J	● — — —		
K	— ● — —		
L	● — ● ●		
M	— —		
N	— ●		
O	— — —		
P	● — — ●		
Q	— — ● —		
R	● — ● ●		
S	● ● ●		
T	—		
		1	● — — — —
		2	● ● — — —
		3	● ● ● — —
		4	● ● ● ● —
		5	● ● ● ● ●
		6	— ● ● ● ●
		7	— — ● ● ●
		8	— — — ● ●
		9	— — — — ●
		0	— — — — —

Relative frequencies  
of letters in the  
English language



# Morse code: Key Idea



A	● —	U	● ● —
B	— ● ● ●	V	● ● ● —
C	— ● — ●	W	● — —
D	— ● ●	X	— ● ● —
E	●	Y	— ● — —
F	● ● — ●	Z	— — ● ●
G	— — ●		
H	● ● ● ●		
I	● ●		
J	● — — —		
K	— ● — —		
L	● — ● ●		
M	— — —		
N	— ●		
O	— — —		
P	● — — ●		
Q	— — ● —		
R	● — ● ●		
S	● ● ●		
T	—		
		1	● — — — —
		2	● ● — — —
		3	● ● ● — —
		4	● ● ● ● —
		5	● ● ● ● ●
		6	— ● ● ● ●
		7	— — ● ● ●
		8	— — — ● ●
		9	— — — — ●
		0	— — — — —

Frequently-used characters are mapped to short codewords.

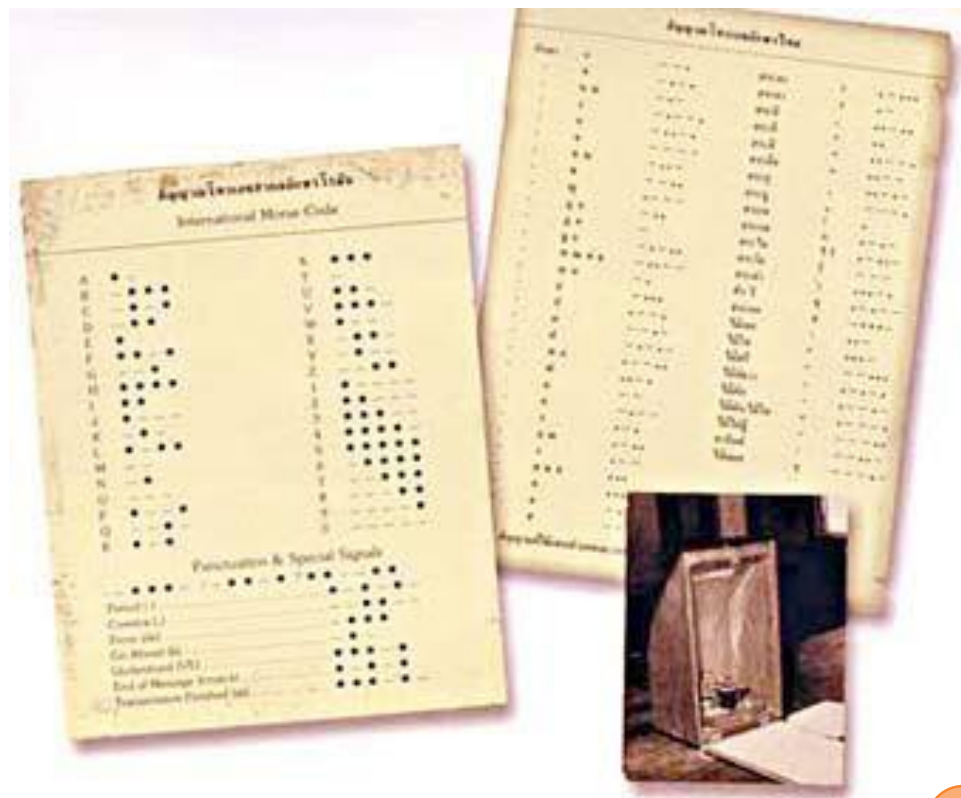


# รหัส Morse ภาษาไทย



รหัสสัญญาณโทรเลขของประเทศไทย  
ที่เริ่มใช้เมื่อ 1 พฤศจิกายน 2455

1. ---	ก	26. ---	ค ข ฅ
2. ---	ข	27. ---	ท
3. ---	ค ฅ	28. ---	ด
4. ---	ด	29. ---	ธ
5. ---	ง	30. ---	ถ
6. ---	จ	31. ---	ช ฌ
7. ---	ช ฌ	32. ---	ซ
8. ---	ซ	33. ---	ญ
9. ---	ญ	34. ---	ฎ
10. ---	ฏ ฎ	35. ---	ฏ
11. ---	ฐ ฎ	36. ---	ฏ
12. ---	ด ฎ	37. ---	ฏ
13. ---	ท ฎ ฅ ฌ	38. ---	ฏ
14. ---	ฌ ฌ	39. ---	ฏ
15. ---	ญ	40. ---	ฏ
16. ---	ฎ	41. ---	ฏ
17. ---	ฎ	42. ---	ฏ
18. ---	ฎ	43. ---	ฎ
19. ---	ท ฎ	44. ---	ฎ
20. ---	ท	45. ---	โทรเลข
21. ---	ท	46. ---	โทรเลข
22. ---	ท	47. ---	การรับ
23. ---	ง	48. ---	โทรเลข
24. ---	ช ฌ	49. ---	โทรเลข
25. ---	ง	50. ---	โทรเลข

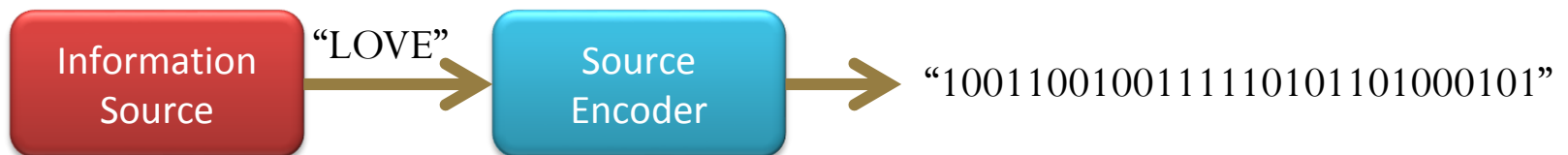


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MATLAB:

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>> X = dec2bin(M, 7);  
>> X = reshape(X', 1, numel(X))  
X =  
1001100100111110101101000101
```



Prof. Robert Mario Fano (MIT)  
Shannon Award (1976 )

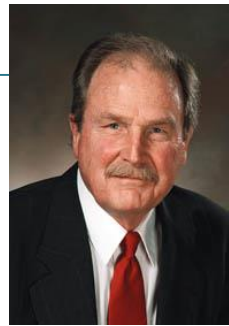


# Shannon–Fano coding

- Proposed in Shannon’s “A Mathematical Theory of Communication” in 1948
- The method was attributed to Fano, who later published it as a technical report.
- Should not be confused with
  - Shannon coding, the coding method used to prove Shannon's noiseless coding theorem, or with
  - Shannon–Fano–Elias coding (also known as Elias coding), the precursor to arithmetic coding.



# Huffman Code



David Huffman (1925–1999)

- MIT, 1951
- Information theory class taught by Professor Fano.
- Huffman and his classmates were given the choice of
  - a term paper on the problem of finding the most efficient binary code.or
  - a final exam.
- Huffman, unable to prove any codes were the most efficient, was about to give up and start studying for the final when he hit upon the idea of using a frequency-sorted binary tree and quickly proved this method the most efficient.
- Huffman avoided the major flaw of the suboptimal Shannon-Fano coding by building the tree from the bottom up instead of from the top down.



# Ex. Huffman Coding in MATLAB

Observe that  
MATLAB  
automatically give  
the **expected**  
**length** of the  
codewords

```
pX = [0.5 0.25 0.125 0.125];           % pmf of X
SX = [1:length(pX)];                   % Source Alphabet
[dict,EL] = huffmandict(SX,pX);       % Create codebook
%% Pretty print the codebook.
codebook = dict;
for i = 1:length(codebook)
    codebook{i,2} = num2str(codebook{i,2});
end
codebook

%% Try to encode some random source string
n = 5; % Number of source symbols to be generated
sourceString = randsrc(1,10,[SX; pX]) % Create data using pX
encodedString = huffmanenco(sourceString,dict) % Encode the data
```



# Ex. Huffman Coding in MATLAB

codebook =

```
[1] '0'  
[2] '1 0'  
[3] '1 1 1'  
[4] '1 1 0'
```

sourceString =

```
1 4 4 1 3 1 1 4 3 4
```

encodedString =

```
0 1 1 0 1 1 0 0 1 1 1 0 0 1 1 0 1 1 1 1 1 0
```





# Ex. Huffman Coding in MATLAB

```
pX = [0.4 0.3 0.1 0.1 0.06 0.04]; % pmf of X
SX = [1:length(pX)]; % Source Alphabet
[dict,EL] = huffmandict(SX,pX); % Create codebook
```

```
%% Pretty print the codebook.
codebook = dict;
for i = 1:length(codebook)
    codebook{i,2} = num2str(codebook{i,2});
end
codebook
```

```
EL
```

The codewords can be different  
from our answers found earlier.  
The expected length is the same.

```
>> Huffman_Demo_Ex2
```

```
codebook =
```

```
 [1]      '1'
 [2]      '0 1'
 [3]      '0 0 0 0'
 [4]      '0 0 1'
 [5]      '0 0 0 1 0'
 [6]      '0 0 0 1 1'
```

```
EL =
```

```
2.2000
```

# Huffman Coding: Source Extension

$X_k \stackrel{\text{i.i.d.}}{\sim} \text{Bernoulli}(p)$   
 $p = 0.1$

